

Recap:

$$H(X|Y) \leq H(X)$$

$$H(X, Y) \leq H(X) + H(Y)$$

Mutual information

$$I(X; Y) = H(X) - H(X|Y) = H(Y) - H(Y|X)$$

KL divergence

$$D(p \parallel q) = \sum_{i=1}^m p_i \log \frac{p_i}{q_i}$$

$$I(X; Y) = D(p(x, y) \parallel p(x) \cdot p(y))$$

Consider a channel with input distribution $p = (p_1, \dots, p_m)$

$$p \rightarrow \boxed{p(y_j | x_i)} \rightarrow r$$

$X \sim$ input distr. channel output distr. $\sim Y$

Let matrix $W = (w_{ij})_{\substack{i=1, \dots, m \\ j=1, \dots, d}} \in \mathbb{R}^{m \times d}$,

W is a stochastic matrix, i.e., $\sum_{j=1}^d w_{ij} = 1, \forall i=1, \dots, m$.

Output distribution $r = (r_1, \dots, r_d)$ is obtained as

$$r = pW$$

Lemma 2.1.13. For any distr. p, q with support $X = \{x_1, \dots, x_m\}$ and stochastic matrix ~~its~~
 $W = (p(y_j | x_i))_{i,j} \in \mathbb{R}^{m \times d}$,
 $D(p \| q) \geq D(pW \| qW)$. \square

Proof. Use the log-sum inequality (a. 2.1.7)

$$\sum_i a_i \log \frac{a_i}{b_i} \geq \sum_i a_i \log \frac{\sum_j a_j}{\sum_j b_j}$$

$$\begin{aligned} D(p \| q) &= \sum_{i=1}^m p(x_i) \log \frac{p(x_i)}{q(x_i)} \\ &= \sum_{i=1}^m \sum_{j=1}^d \underbrace{p(x_i) p(y_j | x_i)}_{a_i} \log \frac{\overbrace{p(x_i) p(y_j | x_i)}^{a_i}}{\underbrace{q(x_i) p(y_j | x_i)}_{b_i}} \\ &\geq \sum_{j=1}^d p \underline{w}_j \log \frac{p \underline{w}_j}{q \underline{w}_j} \\ &= D(pW \| qW) . \quad \square \end{aligned}$$

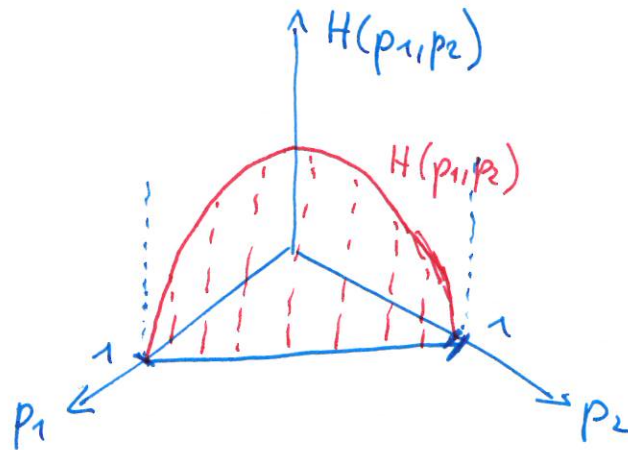
$W = (\underline{w}_1, \dots, \underline{w}_d)$

Theorem 2.1.14.

$H(p)$ is a concave function of $p = (p_1, \dots, p_m)$.

Figure w.r.t. Th. 2.1.14

2-dim. $p = (p_1, p_2)$, $p_1, p_2 \geq 0$, $p_1 + p_2 = 1$



Proof. Let $u = (\frac{1}{m}, \dots, \frac{1}{m})$ be uniform distribution.

$$D(p \| u) = \sum_{i=1}^m p_i \log \frac{p_i}{\frac{1}{m}} = \log m - H(p)$$

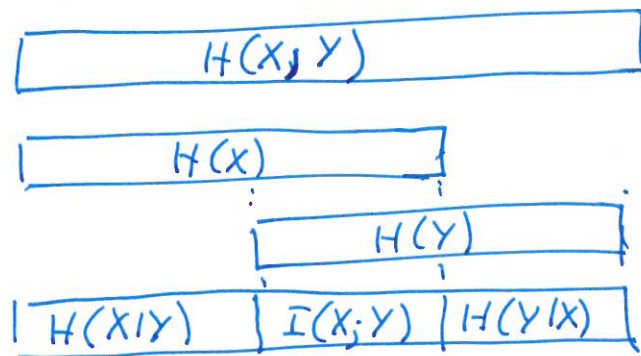
Hence $H(p) = \log m - D(p \| u)$

By Th. 2.1.12 5)

$$\begin{aligned} & D(\lambda p + (1-\lambda)q \| \lambda u + (1-\lambda)u) \\ & \leq \lambda D(p \| u) + (1-\lambda)D(q \| u), \text{ i.e.,} \\ & \text{convexity in the first argument.} \end{aligned}$$

Thus, $H(p)$ is a concave fct. of p . \square

Collecting the quantities in a single picture.



Connections are easy to derive, e.g.,

$$I(X; Y) = H(X) - H(X|Y)$$

$$H(Y|X) = H(X, Y) - H(X)$$

$$I(X; Y) = H(X) + H(Y) - H(X, Y)$$

2.2. Inequalities

Def. 2.2.1. Random variables X, Y, Z are said to have the Markovian property if the joint pmf satisfies

$$p(x, y, z) = p(x) p(y|x) p(z|y).$$

~~For~~ Notation: $X \rightarrow Y \rightarrow Z$.

For $X \rightarrow Y \rightarrow Z$ the conditional distr. of Z depends only on Y and is conditionally independent of X .

Lemma 2.2.2.

- a) If $X \rightarrow Y \rightarrow Z$, then

$$p(x, z | y) = p(x | y) \cdot p(z | y).$$
- b) If $X \rightarrow Y \rightarrow Z$, then $Z \rightarrow Y \rightarrow X$.
- c) If $Z = f(Y)$, then $X \rightarrow Y \rightarrow Z$. \perp

Proof.

a)
$$p(x, z | y) = \frac{p(x, y, z)}{p(y)} = \frac{p(x) p(y | x) p(z | y)}{p(y)}$$

$$= \frac{p(x, y) p(z | y)}{p(y)} = p(x | y) p(z | y).$$

b)
$$p(x, y, z) = p(x) p(y | x) p(z | y)$$

$$= p(x, y) \frac{p(z, y)}{p(y)} \frac{p(z)}{p(z)}$$

$$= p(z) p(y | z) p(x | y), \text{ i.e., } Z \rightarrow Y \rightarrow X.$$

c) $Z = f(Y)$, then

$$p(x, y, z) = \begin{cases} p(x, y), & \text{if } z = f(y) \\ 0, & \text{otherwise} \end{cases}$$

Hence - $p(x, y, z) = p(x, y) \mathbb{1}(z = f(y)) = p(x) p(y | x) p(z | y)$,
 since $p(z | y) = \begin{cases} 1, & z = f(y) \\ 0, & \text{otherwise} \end{cases}$

Theorem 2.2.3. (Data-processing inequality)

If $X \rightarrow Y \rightarrow Z$, then $I(X; Z) \leq \min\{I(X; Y), I(Y; Z)\}$

'No processing of Y can increase the information that Y contains about X .'

Proof By the chain rule ≥ 0

$$I(X; Y, Z) = I(X; Z) + I(X; Y | Z)$$

$$= I(X; Y) + I(X; Z | Y)$$

Since X and Z are cond. independent given Y :

$I(X; Z | Y) = 0$. Since $I(X; Y | Z) \geq 0$, we have

$$I(X; Z) \leq I(X; Y)$$

Equality holds iff $I(X; Y | Z) = 0$, i.e., $X \rightarrow Z \rightarrow Y$.

$I(X; Z) \leq I(Y; Z)$ is shown analogously. \square

Assume X, Y r.v. with support $\mathcal{X} = \{x_1, \dots, x_m\}$.

Define $p_e = P(X \neq Y)$, the 'error probability'.

Theorem 2.2.4. (Fano inequality)

$$H(X|Y) \leq H(p_e) + p_e \log(m-1).$$

This implies that $p_e \geq \frac{H(X|Y) - \log 2}{\log(m-1)}$ \square

$$H(p_e) = -p_e \log p_e - (1-p_e) \log(1-p_e)$$

Beweis Proof.

$$(i) H(X|Y) = \sum_{x \neq y} p(x,y) \log \frac{1}{p(x|y)} + \sum_x p(x,x) \log \frac{1}{p(x|x)}$$

$$(ii) p_e \log(m-1) = \sum_{x \neq y} \log(m-1) p(x,y)$$

$$(iii) H(p_e) = -p_e \log p_e - (1-p_e) \log(1-p_e)$$

$$(iv) \ln t \leq t-1, t \geq 0$$

$$H(X|Y) = p_e \log(m-1) - H(p_e)$$

$$\stackrel{(i),(ii)}{=} \sum_{x \neq y} p(x,y) \log \frac{p_e}{p(x|y)(m-1)}$$

$$+ \sum_x p(x,x) \log \frac{(1-p_e)}{p(x|x)}$$

$$\leq (\log e) \left[\sum_{x \neq y} p(x,y) \left(\frac{p_e}{p(x|y)(m-1)} - 1 \right) \right.$$

$$\left. + \sum_x p(x,x) \left(\frac{1-p_e}{p(x|x)} - 1 \right) \right]$$

$$= (\log e) \left[\frac{p_e}{m-1} \sum_{x \neq y} p(y) - \sum_{x \neq y} p(x,y) \right.$$

$$\left. + (1-p_e) \sum_x p(x) - \sum_x p(x,x) \right]$$

$$= (\log e) [p_e - p_e + (1-p_e) - (1-p_e)] = 0 \quad \blacksquare$$