

Prof. Dr. Rudolf Mathar, Jose Calvo, Markus Rothe

Tutorial 7

- Proposed Solution -

Friday, December 18, 2015

Solution of Problem 1

Algorithm 1 ElGamal signature verification

Require: An ElGamal signature (r, s) , the corresponding message m , a cryptographic hash function h and the corresponding ElGamal public key $y \in Z_p^*$.

Ensure: **True**, if the signature is valid, **False** otherwise

Verify that $1 \leq r \leq p - 1$

$v_1 \leftarrow y^r r^s \pmod{p}$

$v_2 \leftarrow a^{h(m)} \pmod{p}$

if $(v_1 = v_2)$ **then**

return **True**

else

return **False**

end if

1) Verify that $1 \leq r \leq p - 1$, i.e., $1 \leq 373 \leq 848 \checkmark$

2) Compute $v_1 \leftarrow y^r r^s \pmod{p}$:

$$y^r \equiv 399^{373} \equiv 672 \text{ and } r^s \equiv 373^{15} \equiv 643 \pmod{859}.$$

Both results above can be obtained using the Square-and-Multiply algorithm (SQM). Then v_1 yields:

$$v_1 \equiv 672 \cdot 643 \equiv 19 \pmod{859}.$$

3) Compute $v_2 \leftarrow a^{h(m)} \pmod{p}$:

$$v_2 \equiv a^{h(m)} \equiv 206^{65} \equiv 19 \pmod{859}.$$

4) As $v_1 = v_2$ holds, Algorithm 1 returns **True**.

Solution of Problem 2

We have $p \equiv 3 \pmod{4}$, a is a primitive element modulo p , $y \equiv a^x \pmod{p}$, and $a \mid p - 1$. Assume that it is possible to find z such that $a^{rz} \equiv y^r \pmod{p}$, as given in the description. Let $s = \frac{p-3}{2}(m - rz)$.

Task: Show that (r, s) is a valid signature.

Inserting the provided s yields:

$$\begin{aligned} v_1 &\equiv y^r r^s \equiv a^{rz} r^{\frac{p-3}{2}(m-rz)} \\ &\equiv a^{rz} (r^{\frac{p-3}{2}})^{m-rz} \pmod{p}. \end{aligned} \quad (1)$$

From $a \mid p - 1$ it follows that there exists a $v \in \mathbb{Z}$ such that $va = p - 1$.

Now, choose $r = v$:

$$ra \equiv p - 1 \pmod{p} \Leftrightarrow r \equiv a^{-1}(p - 1) \equiv -(a^{-1}) \pmod{p}.$$

To obtain (1), we first substitute r and exponentiate it by the power of $\frac{p-3}{2}$:

$$\Leftrightarrow r^{\frac{p-3}{2}} \equiv (-(a^{-1}))^{\frac{p-3}{2}} \pmod{p}.$$

Note that $(-1) \pmod{p}$ is self-inverse:

$$\Leftrightarrow r^{\frac{p-3}{2}} \equiv \left((-a)^{\frac{p-3}{2}} \right)^{-1} \pmod{p}.$$

For $\frac{p-3}{2}$ even, we obtain $(-1)^{\frac{p-3}{2}} = 1$, and with that:

$$\begin{aligned} &\Rightarrow r^{\frac{p-3}{2}} \equiv \left((-1)^{\frac{p-3}{2}} a^{\frac{p-3}{2}} \right)^{-1} \\ &\equiv (a^{\frac{p-3}{2}})^{-1} \equiv a^{-\frac{p-3}{2}} \\ &\equiv a^{-(\frac{p-1}{2}-1)} \equiv a^{-\frac{p-1}{2}+1} \\ &\equiv \underbrace{a^{-\frac{p-1}{2}}}_{\equiv -1 \pmod{p}} a \equiv -a \pmod{p}. \end{aligned}$$

For the last line, note that a is a primitive element and that $\left(a^{\frac{p-1}{2}}\right)^2 \equiv 1 \pmod{p}$.

This result provides the following for (1):

$$\begin{aligned} v_1 &\equiv y^r r^s \equiv a^{rz} r^{\frac{p-3}{2}(m-rz)} \\ &\equiv a^{rz} (-a)^{(m-rz)} \\ &\equiv a^{rz} a^{m-rz} (-1)^{(m-rz)} \pmod{p}. \end{aligned}$$

Choose m such that $m - rz$ is even valued:

$$\begin{aligned} v_1 &\equiv a^{rz} a^{(m-rz)} \\ &\equiv a^m \equiv v_2 \pmod{p}, \end{aligned}$$

so that the forged signature is valid.

Solution of Problem 3

In the ElGamal verification $v_1 \equiv v_2 \pmod{p}$ needs to be fulfilled.

Recall that $y = a^x \pmod{p}$ and $r = a^k \pmod{p}$ are used:

$$\begin{aligned} y^r r^s &\equiv a^{h(m)} \pmod{p} \\ \Leftrightarrow a^{xr} a^{ks} &\equiv a^{h(m)} \pmod{p} \\ \stackrel{\text{Fermat}}{\Leftrightarrow} xr + ks &\equiv h(m) \pmod{p-1}. \end{aligned}$$

Now, we expand both sides of the congruence with $h(m)^{-1}h(m')$:

$$xr \cdot h(m)^{-1}h(m') + ks \cdot h(m)^{-1}h(m') \equiv h(m)h(m)^{-1}h(m') \equiv h(m') \pmod{p-1} \quad (2)$$

$$\Leftrightarrow xr' + ks' \equiv h(m') \pmod{p-1} \quad (3)$$

$$\begin{aligned} \stackrel{\text{Fermat}}{\Leftrightarrow} a^{xr'} a^{ks'} &\equiv a^{h(m')} \pmod{p} \\ \Leftrightarrow y^{r'} r^{s'} &\equiv a^{h(m')} \pmod{p} \\ \stackrel{\downarrow}{\Leftrightarrow} y^{r'} (r')^{s'} &\equiv a^{h(m')} \pmod{p}. \end{aligned}$$

The equivalence assumption in the last line holds if $r \equiv r' \pmod{p}$.

Note: In the ElGamal scheme, the condition $1 \leq r < p$ must be checked!

From (2) and (3), we have $rh(m)^{-1}h(m') \equiv r' \pmod{p-1}$.

We have to solve the following system of two congruences w.r.t. r' :

$$\begin{aligned} r' &\equiv rh(m)^{-1}h(m) \pmod{p-1}, \\ r' &\equiv r \pmod{p}. \end{aligned}$$

By means of the Chinese Remainder Theorem, we get the parameters:

$$\begin{aligned} a_1 &= r \pmod{p}, & a_2 &= rh(m)^{-1}h(m') \pmod{p-1}, \\ m_1 &= p, & m_2 &= p-1, \\ M_1 &= p-1, & M_2 &= p, \\ y_1 &= M_1^{-1} \equiv p-1 \pmod{p}, & y_2 &= M_2^{-1} \equiv 1 \pmod{p-1}, \\ M &= p(p-1). \end{aligned}$$

The Chinese Remainder Theorem leads to the solution:

$$\begin{aligned} r' &= \sum_{i=1}^2 a_i M_i y_i = r(p-1)^2 + rh(m)^{-1}h(m')p \\ &\equiv r(p^2 - p - p + 1 + h(m)^{-1}h(m')p) \\ &\equiv r(p(p-1) - p + 1 + h(m)^{-1}h(m')p) \\ &\equiv r(h(m)^{-1}h(m')p - p + 1) \pmod{M}. \end{aligned}$$

The forged signature

$$(r', s') = (r(h(m)^{-1}h(m')p - p + 1) \pmod{M}, sh(m)^{-1}h(m') \pmod{(p-1)})$$

is a valid signature of $h(m')$, if $1 \leq r < p$ is not checked.

Solution of Problem 4

Recall for a), b) and c) that we have: $r = a^k \pmod{p}$ and $y = a^x \pmod{p}$ from the ElGamal signature scheme.

- a) This is easily solved by substituting $s = x^{-1}(h(m) - kr)$, r and y :

$$\begin{aligned} v_1 &\equiv y^s r^r \equiv y^{x^{-1}(h(m)-kr)} a^{kr} \\ &\equiv a^{x x^{-1}(h(m)-kr)} a^{kr} \\ &\equiv a^{(h(m)-kr)+kr} \\ &\equiv a^{h(m)} \equiv v_2 \pmod{p}. \end{aligned}$$

If the given signature is properly checked, $v_1 = y^s r^r = a^{h(m)} = v_2 \pmod{p}$ is true.

- b) In this case it is useful to proceed stepwise. We begin with computing:

$$a^s \equiv a^{xh(m)+kr} \equiv a^{xh(m)} a^{kr} \pmod{p}.$$

Next, we substitute y and r , correspondingly, and we rearrange the congruence:

$$\begin{aligned} a^s &\equiv y^{h(m)} r^r \pmod{p} \\ \Leftrightarrow a^s r^{-r} &\equiv y^{h(m)} \pmod{p}. \end{aligned}$$

In the last step, we fix the parameters for verification by:

$$\begin{aligned} v_1 &\equiv a^s r^{-r} \pmod{p}, \\ v_2 &\equiv y^{h(m)} \pmod{p}, \end{aligned}$$

so that $v_1 = v_2$ must be checked by the proposed scheme.

- c) In analogy to b), we compute:

$$\begin{aligned} a^s &\equiv a^{xr+kh(m)} \\ &\equiv a^{xr} a^{kh(m)} \\ &\equiv y^r r^{h(m)} \pmod{p} \\ \Leftrightarrow v_1 &= a^s y^{-r} \equiv r^{h(m)} = v_2 \pmod{p}. \end{aligned}$$