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Tutorial 6

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Problem 1. (*Blum-Blum-Shub generator*) The security of the Blum-Blum-Shub generator is based on the difficulty to compute square roots modulo $n = pq$ for two distinct primes p and q with $p, q \equiv 3 \pmod{4}$.

Design a generator for pseudo-random bits which is based on the hardness of the RSA-problem.

Problem 2. (*number of messages and hardware resources of two hash functions*) Consider two hash functions, one with an output length of 64 bits and another one with an output length of 128 bits.

For each of these functions, do the following:

- a) Determine the number of messages that have to be created to find a collision with a probability larger than 0.86 by means of the birthday paradox.
- b) Determine the hardware resources required for this attack in terms of memory size, number of comparisons, and number of hash function executions.

Problem 3. (*CBC and CFB for MAC generation*) Both, the CBC mode and the CFB mode, can be used for the generation of a MAC as follows.

- A plaintext is divided into n equally-sized blocks M_1, \dots, M_n .
- For the CFB-MAC, the ciphertexts are $C_i = M_{i+1} \oplus E_K(C_{i-1})$ for $i = 1, \dots, n-1$ and $\text{MAC}_K^{(n)} = E_K(C_{n-1})$ with initial value $C_0 = M_1$.
- For the CBC-MAC, the ciphertexts are $\hat{C}_i = E_K(\hat{C}_{i-1} \oplus M_i)$ for $i = 1, \dots, n-1$ and $\widehat{\text{MAC}}_K^{(n)} = E_K(\hat{C}_{n-1} \oplus M_n)$ with initial value $\hat{C}_0 = 0$.

Show that the equivalency $\text{MAC}_K^{(n)} = \widehat{\text{MAC}}_K^{(n)}$ holds.

Problem 4. (*derive a message validation protocol*) Suppose Alice transmits the following cryptogram to Bob:

$$c = e(m \parallel h(k_2 \parallel m), k_1).$$

Assume that the message m , the shared keys k_1, k_2 , the hash values $h(x)$ and the output of the encryption function have fixed lengths known to Alice and Bob.

- a) Derive a protocol for decryption and message validation used by Bob.
- b) Modify the given scheme to construct a similar protocol for a public-key cryptosystem. You may use two private-/public key-pairs (K_1, L_1) and (K_2, L_2) and a session key s used in the hash, which is securely transmitted to Bob within the cryptogram c .
- c) How can an intruder Eve impersonate Alice to Bob in the system of (b)? How could the attack be prevented?