

Review Exercise

Advanced Methods of Cryptography

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Problem 4.

Alice and Bob use the ElGamal signature scheme to sign messages. In order to reduce the amount of computation, Alice signs $h(m)$ (instead of signing a message m directly) with h as hash function given by

$$h(m) = m(m + 3) \bmod n,$$

with $n = uv$ and u, v are prime numbers.

- List the four requirements that the function h should fulfill to be used as cryptographic hash function.
- Is h collision free in general?

Alice wants to sign $h(m)$ with $m = 83$ using the ElGamal signature scheme. She chooses the prime number $p = 101$ and the parameter $a = 7$.

- Calculate the hash value $h(m)$ given $u = 3$ and $v = 19$.
- Which condition must be fulfilled by a , to be used in the ElGamal signature scheme? Show that $a = 7$ fulfills this condition.
- Alice chooses the private key $x_A = 11$ and the random secret $k = 17$. Compute the signature (r, s) of $h(m)$.
- Bob receives a signature $(r, s) = (120, 67)$ (this is not the signature from (e)). Is this signature valid?

Problem 5.

In the following, a certification and identification protocol is considered. It is an extension of the Fiat-Feige-Shamir protocol. It establishes authentication between A and B with the aid of a trusted authority server T . The utilized parameters are denoted as a public $n = pq$ of two secret, large primes p, q with $p \neq q$, A 's private key $u \in \mathbb{Z}_n^*$, A 's public key $v \in \mathbb{Z}_n^*$, a random number $r \in \mathbb{Z}_n \setminus \{p, q\}$, a random number c and a large publicly known exponent $e \in \mathbb{Z}_{\varphi(n)}$. Furthermore, a signature algorithm S_T used by T , a verification algorithm V_T , a signature s and a public certificate $cert_T(A)$ issued by T to A are used.

Certification and Identification Protocol

- (1) A computes $v = (u^{-1})^e \pmod{n}$.
 $A \rightarrow T : v$
- (2) T computes $s = S_T(A, v)$ and $\text{cert}_T(A) = (A, v, s)$.
 $T \rightarrow A : \text{cert}_T(A)$
- (3) A chooses a random $r \in \mathbb{Z}_n \setminus \{p, q\}$ and computes $x = r^e \pmod{n}$.
 $A \rightarrow B : x, \text{cert}_T(A)$
- (4) B checks $V_T(s)$ and matches (A, v) from $\text{cert}_T(A)$.
If both are valid, B chooses a random $c \in \{1, \dots, e\}$.
 $B \rightarrow A : c$
- (5) A computes $y = ru^c \pmod{n}$.
 $A \rightarrow B : y$
- (6) B verifies $x \equiv y^e v^c \pmod{n}$.

Answer the following questions with respect to this protocol.

- (a) What is the purpose of the certificate $\text{cert}_T(A)$ in this protocol?
- (b) Name the two number-theoretic problems this protocol relies on.
- (c) Prove that the verification works.
- (d) The security of this protocol also relies on cryptographically secure random numbers. Calculate the random sequence b_1, \dots, b_t with the Blum-Blum-Shub Generator using $x_0 = 29$, $n = 13 \cdot 23$ and $t = 5$. The given numbers are decimal. Is $x_0 = 29$ a valid initial value? Reason your statement.

Problem 6.

Consider the equation

$$Y^2 = X^3 + X + 3.$$

- (a) Show that this equation describes an elliptic curve E over the field \mathbb{F}_7 .
- (b) Calculate all points on $E(\mathbb{F}_7)$. What is the order of $E(\mathbb{F}_7)$?
- (c) For each point on $E(\mathbb{F}_7)$, calculate its inverse.
- (d) For each point on $E(\mathbb{F}_7)$, calculate its order.
- (e) Is the group $E(\mathbb{F}_7)$ cyclic?
- (f) Find all solutions of the equation $4P = \mathcal{O}$ in $E(\mathbb{F}_7)$.