

## Homework 7 in Cryptography II

Prof. Dr. Rudolf Mathar, Wolfgang Meyer zu Bergsten, Michael Reyer  
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**Exercise 20.** Consider two hash functions, one with an output length of 64 bits and another one with an output length of 128 bits.

For each of these functions, do the following:

- Determine the number of messages that have to be created to find a collision with a probability larger than 0.86 by means of the birthday paradox.
- Determine the hardware resources required for this attack in terms of memory size, number of comparisons and number of hash function executions.

**Exercise 21.** Using a block cipher  $E_K(x)$  with block length  $k$  and key  $K$  a hash function  $h(m)$  is provided in the following way:

Append  $m$  with zero bits until it is a multiple of  $k$ , divide  $m$  into  $n$  blocks of  $k$  bits.

$c \leftarrow E_{m_0}(m_0)$

**for**  $i$  **in**  $1 \dots (n - 1)$

$d \leftarrow E_{m_0}(m_i)$

$c \leftarrow c \oplus d$

**end for**

$h(m) \leftarrow c$

Does this function fulfill the basic requirements for a cryptographic hash function? Can these requirements be fulfilled by replacing the XOR-operation by a logical AND?

**Exercise 22.** Besides the CBC mode, the CFB mode can be used for the generation of a MAC. The plaintext consists of the blocks  $M_1, \dots, M_n$ , and we set the initialization vector  $C_0 := M_1$ . Now, we encrypt  $M_2, \dots, M_n$  in CFB mode with key  $K$ , which results in the ciphertexts  $C_1, \dots, C_{n-1}$ . For the MAC, we use  $MAC_K := E_K(C_{n-1})$ .

Show that this scheme results in the same MAC as the algorithm in example 10.5 from the lecture notes with the initial value set to  $C_0 := \mathbf{0}$ .